

A CATEGORICAL OUTLOOK ON CELLULAR AUTOMATA

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ABSTRACT. In programming language semantics, it has proved to be fruitful to analyze context-dependent notions of computation, e.g., dataflow computation and attribute grammars, using comonads. We explore the viability and value of similar modeling of cellular automata. We identify local behaviors of cellular automata with coKleisli maps of the exponent comonad on the category of uniform spaces and uniformly continuous functions and exploit this equivalence to conclude some standard results about cellular automata as instances of basic category-theoretic generalities. In particular, we recover Ceccherini-Silberstein and Coornaert’s version of the Curtis-Hedlund theorem.

1. Introduction

Since the seminal work of Moggi [13], it has become standard in programming language semantics to analyze functions producing effects such as exceptions, input, output, interactive input-output, nondeterminism, probabilistic choice, state, continuations using monads. Specifically, effectful functions are identified with Kleisli maps of a suitable monad on the category of pure functions. Wadler [18] put this view to further use in programming methodology by extracting from it a very useful programming idiom for purely functional languages like Haskell.

The dual view of context-dependent functions as coKleisli maps of a comonad is equally useful, but less well known. Brookes and Geva [2] explained the “intensional” aspect of denotational semantics in terms of the ω -chain comonad on the category of ω -cpos. More recently, Uustalu and Vene [15, 16, 17] employed comonads to analyze dataflow computation and attribute grammars and Hasuo et al. [9] treated tree transducers.

Characteristic of many context-dependent notions of computation is shape-preserving transformation of some datastructure based on a value update rule which is local in nature and applied uniformly to every node. This is the case with dataflow computation where such a transformation is applied to a list or to a stream with a distinguished position and with attribute grammars where computation happens on a tree or a tree with a distinguished position (a “zipper”). Cellular automata are similar, the datastructure being the integer line or plane or, more generally, a group.

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It should therefore be worthwhile to test the slogan that context-dependent computation is comonadic also on cellular automata. To a degree, this has already been done, as Piponi [14] programmed cellular automata in Haskell using the comonadic interface. However, he did not use his modeling of cellular automata to prove properties about them and also dropped the classical requirement that cellular automata rely on a finite neighborhood only.

In this paper, we study the comonadic aspect of cellular automata deeper. We identify cellular automata (more exactly their local behaviors) with coKleisli maps of the exponent monad on **Unif**, the category of uniform spaces and uniformly continuous functions, and explore whether this view can be useful. We see that it is: we can conclude some standard results about cellular automata as instances of category-theoretic generalities. In particular, we recover the Curtis-Hedlund theorem [10]—a characterization of global behaviors of cellular automata—in the version of Ceccherini-Silberstein and Coornaert [4] (this applies to general discrete alphabets rather than finite alphabets only). This theorem turns out to be an instance of the basic category-theoretic fact that the coKleisli category of a comonad is isomorphic to the full subcategory of its co-Eilenberg-Moore category given by the cofree coalgebras. We also show that the comonadic view allows one to see 2-dimensional cellular automata as 1-dimensional and treat point-dependent cellular automata.

The paper is organized as follows. Section 2 is a quick introduction to comonads while Section 3 reviews some preliminaries about topological and uniform spaces. In Section 4, we show that cellular automata local behaviors are the same as coKleisli maps of a certain comonad. In Section 5, we recover the Curtis-Hedlund theorem (in the version of Ceccherini-Silberstein and Coornaert). In Section 6, we reprove the reversibility principle. In Sections 7, 8, we discuss some further applications of the comonadic view: 2-dimensional cellular automata as 1-dimensional and point-dependent cellular automata.

The paper assumes knowledge of basic category theory (categories, functors, natural transformations, Cartesian closed categories), but is self-contained in regards to comonads. For background material on category theory and (co)monads in particular, we refer the reader to Barr and Wells [1, Ch. 1, 3]. We also assume the basics of CA as presented by Ceccherini-Silberstein and Coornaert [5, Ch. 1].

2. Comonads

Given two categories \mathcal{C} , \mathcal{D} and a functor $L : \mathcal{D} \rightarrow \mathcal{C}$, a *right adjoint* to L is given by a functor $R : \mathcal{C} \rightarrow \mathcal{D}$ and two natural transformations $\varepsilon : LR \rightarrow \text{Id}_{\mathcal{C}}$ (the *counit*) and $\eta : \text{Id}_{\mathcal{D}} \rightarrow RL$ (the *unit*) such that the diagrams

$$\begin{array}{ccc} L & \xrightarrow{L\eta} & LRL \\ & \searrow & \downarrow \varepsilon L \\ & & L \end{array} \quad \begin{array}{ccc} R & & \\ \eta R \downarrow & \searrow & \\ RLR & \xrightarrow{R\varepsilon} & R \end{array}$$

commute. Equivalently, a right adjoint may be given by an object mapping $R : |\mathcal{C}| \rightarrow |\mathcal{D}|$, for any object $A \in |\mathcal{C}|$, a map $\varepsilon_A : LRA \rightarrow A$, and, for any objects $A \in |\mathcal{D}|$, $B \in |\mathcal{C}|$ and map $k : LA \rightarrow B$, a map $k^\ddagger : A \rightarrow RB$ (the *right transpose*) such that

- for any objects $A \in |\mathcal{D}|$, $B \in |\mathcal{C}|$, and map $k : LA \rightarrow B$, $\varepsilon_B \circ Lk^\ddagger = k$,

- for any object $A \in |\mathcal{C}|$, $(\varepsilon_A)^\ddagger = \text{id}_{RA}$,
- for any objects $A, B \in |\mathcal{D}|$, $C \in |\mathcal{C}|$ and maps $f : A \rightarrow B$, $k : LB \rightarrow C$, $(k \circ Lf)^\ddagger = k^\ddagger \circ f$.

The morphism mapping part of R and unit η define the right transpose of $k : LA \rightarrow B$ by $k^\ddagger =_{\text{df}} Rk \circ \eta_A$. The right transpose $(-)^{\ddagger}$ determines the morphism mapping part of R and unit η by $Rf =_{\text{df}} (f \circ \varepsilon_A)^\ddagger$, for $f : A \rightarrow B$, and $\eta_A =_{\text{df}} (\text{id}_{LA})^\ddagger$.

A *comonad* on a category \mathcal{C} is given by a functor $D : \mathcal{C} \rightarrow \mathcal{C}$ and natural transformations $\varepsilon : D \rightarrow \text{Id}_{\mathcal{C}}$ (the *counit*) and $\delta : D \rightarrow DD$ (the *comultiplication*) making the diagrams

$$\begin{array}{ccc} D & \xrightarrow{\delta} & DD \\ \delta \downarrow & \searrow & \downarrow \varepsilon D \\ DD & \xrightarrow{D\varepsilon} & D \end{array} \quad \begin{array}{ccc} D & \xrightarrow{\delta} & DD \\ \delta \downarrow & & \downarrow \delta D \\ DD & \xrightarrow{D\delta} & DDD \end{array}$$

commute. Equivalently, a comonad can be given by an object mapping $D : |\mathcal{C}| \rightarrow |\mathcal{C}|$, for any object $A \in |\mathcal{C}|$, a map $\varepsilon_A : DA \rightarrow A$, and, for any objects $A, B \in |\mathcal{C}|$ and map $k : DA \rightarrow B$, a map $k^\ddagger : DA \rightarrow DB$ (the *coKleisli extension*) such that

- for any objects $A, B \in |\mathcal{C}|$ and map $k : DA \rightarrow B$, $\varepsilon_B \circ k^\ddagger = k$,
- for any object A , $(\varepsilon_A)^\ddagger = \text{id}_{DA}$,
- for any objects $A, B, C \in |\mathcal{C}|$ and maps $k : DA \rightarrow B$, $\ell : DB \rightarrow C$, $(\ell \circ k^\ddagger)^\ddagger = \ell^\ddagger \circ k^\ddagger$.

The morphism mapping part of D and comultiplication δ define the coKleisli extension $(-)^{\ddagger}$ by $k^\ddagger =_{\text{df}} Dk \circ \delta_A$. Conversely, the $(-)^{\ddagger}$ determines the morphism mapping part of D and comultiplication δ by $Df =_{\text{df}} (f \circ \varepsilon_A)^\ddagger$, $\delta_A =_{\text{df}} (\text{id}_{DA})^\ddagger$.

A functor $L : \mathcal{C} \rightarrow \mathcal{D}$ with a right adjoint (R, ε, η) , defines a comonad on \mathcal{C} with counit ε by $D =_{\text{df}} LR$, $\delta =_{\text{df}} L\eta R$, alternatively by $DA =_{\text{df}} L(RA)$, $k^\ddagger =_{\text{df}} Lk^\ddagger$.

In the converse direction, a comonad (D, ε, δ) on \mathcal{C} induces a whole category of adjunctions $(\mathcal{D}, L, R, \eta)$ that have ε as the counit and satisfy $D = LR$, $\delta = L\eta R$, called *splittings* of the comonad. This category has initial and final objects, which are known as the coKleisli and coEilenberg-Moore splittings of the comonad.

The *coKleisli category* $\text{coKl}(D)$ has as objects those of \mathcal{C} and as maps from A to B those from DA to B of \mathcal{C} . The identity jd_A on object A is defined by $\text{jd}_A =_{\text{df}} \varepsilon_A$. The composition $\ell \bullet k$ of maps $k : DA \rightarrow B$ and $\ell : DB \rightarrow C$ is $\ell \bullet k =_{\text{df}} \ell \circ Dk \circ \delta_A = \ell \circ k^\ddagger$. The functor $L : \text{coKl}(D) \rightarrow \mathcal{C}$ in the coKleisli splitting is defined by $LA =_{\text{df}} DA$, $Lk =_{\text{df}} k^\ddagger$. The right adjoint, unit and right transpose are defined by $RA =_{\text{df}} A$, $Rf =_{\text{df}} f \circ \varepsilon_A$, $\eta_A =_{\text{df}} \text{id}_{DA}$, $k^\ddagger =_{\text{df}} k$.

The *coEilenberg-Moore category* $\text{coEM}(D)$ has as objects coalgebras of D and as maps coalgebra maps of D . A *coalgebra* of D is given by an object $A \in |\mathcal{C}|$ and map $u : A \rightarrow DA$ (the *coalgebra structure*) making the diagrams

$$\begin{array}{ccc} A & \xrightarrow{u} & DA \\ & \searrow & \downarrow \varepsilon_A \\ & & A \end{array} \quad \begin{array}{ccc} A & \xrightarrow{u} & DA \\ u \downarrow & & \downarrow \delta_A \\ DA & \xrightarrow{Du} & D(DA) \end{array}$$

commute. A *coalgebra map* between (A, u) and (B, v) is a map $f : A \rightarrow B$ making the diagram

$$\begin{array}{ccc} A & \xrightarrow{u} & DA \\ f \downarrow & & \downarrow Df \\ B & \xrightarrow{v} & DB \end{array}$$

commute. The identity and composition are inherited from \mathcal{C} . The functor $L : \text{coEM}(D) \rightarrow \mathcal{C}$ in the splitting of D through $\text{coEM}(D)$ is the coalgebra-structure forgetful functor: $L(A, u) =_{\text{df}} A$, $Lf =_{\text{df}} f$. The right adjoint R is defined by $RA =_{\text{df}} (DA, \delta_A)$, $Rf =_{\text{df}} Df$. The unit and right transpose are defined by $\eta_{(A, u)} =_{\text{df}} u$ and $k^\ddagger =_{\text{df}} Dk \circ u$ (for $k : L(A, u) \rightarrow B$).

The functor R being the right adjoint of the forgetful functor implies that, for any B , the coalgebra $RB = (DB, \delta_B)$ is the cofree coalgebra on B , i.e., for any coalgebra (A, u) , object B and map $k : A \rightarrow B$, there is a unique map $f : A \rightarrow DB$, namely k^\ddagger , such that the diagrams

$$\begin{array}{ccccc} & & A & \xrightarrow{u} & DA \\ & & \downarrow f & & \downarrow Df \\ k \swarrow & & & & \\ B & \xleftarrow{\varepsilon_B} & DB & \xrightarrow{\delta_B} & D(DB) \end{array}$$

commute.

The unique splitting map between the coKleisli and coEilenberg-Moore splitting is the functor $E : \text{coKl}(D) \rightarrow \text{coEM}(D)$ defined by $EA =_{\text{df}} (DA, \delta_A)$, $Ek =_{\text{df}} Dk \circ \delta_A = k^\ddagger$. This functor is a full embedding. The image of E is the full subcategory of $\text{coEM}(D)$ given by the cofree coalgebras that is therefore isomorphic to $\text{coKl}(D)$.

A simple and instructive example of a comonad and its coKleisli and coEilenberg-Moore splittings is given by the *reader* (or *product*) *comonad*. It is defined on any category \mathcal{C} with finite products, but let us choose \mathcal{C} to be **Set** (or **Top** or **Unif**), so we can write pointwise definitions for intuitiveness. Given some fixed object $C \in |\mathcal{C}|$, it is defined by $DA =_{\text{df}} A \times C$, $Df(x, c) =_{\text{df}} (f(x), c)$, $\varepsilon_A(x, c) =_{\text{df}} x$, $\delta_A(x, c) =_{\text{df}} ((x, c), c)$, $k^\ddagger(x, c) =_{\text{df}} (k(x, c), c)$.

The coKleisli category has as objects those of \mathcal{C} and as maps from A to B those from $A \times C$ to B . The identities and composition are defined by $\text{jd}(x, c) =_{\text{df}} x$, $(\ell \bullet k)(x, c) =_{\text{df}} \ell(k(x, c), c)$.

A coalgebra of D is given by an object A and a map $u : A \rightarrow A \times C$ satisfying the laws of a coalgebra. Let us define $(u_0(x), u_1(x)) =_{\text{df}} u(x)$. The laws impose that $u_0(x) = x$ and $((u_0(x), u_1(x)), u_1(x)) = (u(u_0(x)), u_1(x))$. The first law defines u_0 and the second becomes a tautology as soon as this definition is substituted into it. Hence, a coalgebra is effectively the same as an object A with an unconstrained map $u_1 : A \rightarrow C$.

A map between D -coalgebras (A, u) , (B, v) is a map $f : A \rightarrow B$ such that $(f(u_0(x)), u_1(x)) = (v_0(f(x)), v_1(f(x)))$, which boils down to $u_1(x) = v_1(f(x))$.

The coEilenberg-Moore category has thus as objects pairs of an object A and map $u_1 : A \rightarrow C$ and a map between (A, u_1) , (B, v_1) is map $f : A \rightarrow B$ such that $u_1(x) = v_1(f(x))$. The cofree coalgebra on A is the pair $(A \times C, \delta_{1A})$ where $\delta_{1A}(x, c) =_{\text{df}} c$.

The isomorphism between $\text{coKl}(D)$ and the category of cofree D -coalgebras establishes a 1-1 correspondence between maps $k : A \times C \rightarrow B$ and maps $f : A \times C \rightarrow B \times C$ such that $\delta_{1B}(f(x, c)) = c$.

3. Exponentials, topologies, and uniformities

Given an object C in a category \mathcal{C} with finite products, it is said to be *exponentiable* if the functor $(-)\times C$ has a right adjoint. This amounts to the existence, for any object A , of an object A^C (the *exponential*) and map $\text{ev}_A : A^C \times A \rightarrow C$ (the *evaluation*) as well as, for any objects A, B and map $k : A \times C \rightarrow B$, a map $\text{cur}(k) : A \rightarrow B^C$ (the *currying* of k) satisfying appropriate conditions. If every object of \mathcal{C} is exponentiable, it is called *Cartesian closed*. Intuitively, exponentials are internalized homsets. In **Set**, every object C is exponentiable and the exponential A^C is the set of all functions from C to A .

Things are somewhat more complicated in the category **Top** of topological spaces, as the exponential A^C is to be the set of continuous functions from C to A , but it must also be given a topology. Moreover, the evaluation $\text{ev}_A : A^C \times A \rightarrow C$ must be continuous and the currying $\text{cur}(k) : A \rightarrow B^C$ of a continuous function must be continuous.

Not every topological space is exponentiable. Hausdorff spaces are exponentiable if and only if they are locally compact: in this case, the exponential topology on the space of continuous functions from C to A is the *compact-open topology* generated by the sets $\{f : C \rightarrow A \mid f(K) \subseteq U\}$ with K compact in C and U open in A [7, 8]. In particular, discrete spaces are exponentiable (which also follows from the discrete topology making every function from it continuous) and their compact-open topology is in fact the product topology.

That not all objects can act as exponents is also true in the category **Unif** of uniform spaces whose constituents we now define. A *uniform space* is a set A endowed with a *uniformity*, i.e., a collection \mathcal{U} of binary relations on A (called *entourages*) satisfying the following properties:

- (1) $\Delta \subseteq U$ for every $U \in \mathcal{U}$, where $\Delta = \{(x, x) \mid x \in A\}$ is the *diagonal*.
- (2) If $U \subseteq V$ and $U \in \mathcal{U}$ then $V \in \mathcal{U}$.
- (3) If $U, V \in \mathcal{U}$ then $U \cap V \in \mathcal{U}$.
- (4) If $U \in \mathcal{U}$ then $U^{-1} \in \mathcal{U}$.
- (5) If $U \in \mathcal{U}$ then $V^2 = \{(x, y) \mid \exists z \mid (x, z), (z, y) \in V\} \subseteq U$ for some $V \in \mathcal{U}$.

The simplest non-trivial uniformity on A is the *discrete uniformity*, made of all the supersets of the diagonal. A uniformity induces a topology as follows: $\Omega \subseteq A$ is open if and only if, for every $x \in \Omega$, there exists $U \in \mathcal{U}$ such that $\{y \in A \mid (x, y) \in U\} \subseteq \Omega$. Such topology is Hausdorff if and only if $\bigcap_{U \in \mathcal{U}} U = \Delta$. The discrete uniformity induces the discrete topology, but is not the only one that does (cf. [11, I-5]), i.e., uniform spaces may be discrete without being *uniformly discrete*.

A map $f : A \rightarrow B$ between uniform spaces is *uniformly continuous* (briefly, u.c.) if, for every entourage V on B , there is an entourage U on A such that $(f \times f)(U) \subseteq V$. Any u.c. function is continuous in the topology induced by the uniformities: the converse is true if A is compact [11, II-24] but false in general even for metric spaces. The *product uniformity* is the coarsest uniformity that makes the projections uniformly continuous: the topology induced by the product uniformity is the product topology. A product of discrete uniformities is called *prodiscrete*.

In **Unif**, uniformly discrete objects are exponentiable [11, III.19 and III.21]. Again, the reason is that every function from C is u.c. as soon as C is uniformly discrete.

4. Cellular automata as coKleisli maps

Classically, a *cellular automaton* on a monoid $(G, 1_G, \cdot)$ (the *universe*)¹ and set A (the *alphabet*)² is given by a finite subset N of G (the *support neighborhood*) and function $d : A^N \rightarrow A$ (the *transition rule*).

Any cellular automaton induces a *local behavior* $k : A^G \rightarrow A$ via $k(c) =_{\text{df}} d(c|_N)$.

One speaks of elements A^N of finite subsets $N \subseteq G$ as *patterns* and elements of A^G as *configurations*. Transition rules work on patterns, local behaviors on configurations. Cellular automata that induce the same local behavior are considered equivalent. In this paper, we will not distinguish between equivalent cellular automata, hence we can identify cellular automata with their local behaviors.

If the alphabet A is finite, a function $k : A^G \rightarrow A$ is a local behavior (i.e., the local behavior of some cellular automaton) if and only if it is continuous for the discrete topology on A and the product topology on A^G .

In the general case (where A may be infinite), the above equivalence does not generally hold, but a refinement does. A function $k : A^G \rightarrow A$ is then a local behavior iff it is uniformly continuous for the discrete uniformity on A and product uniformity on A^G .³ The finite case becomes an instance: if A is finite, then A^G is compact and therefore any continuous function $k : A^G \rightarrow A$ is uniformly continuous.

Based on these observations, we henceforth take it as a definition that a local behavior on a set (the alphabet) A is a uniformly continuous function $k : A^G \rightarrow A$ wrt. the prodiscrete uniformity on A^G and forget about the definition of cellular automata in terms of a support neighborhood and a transition rule.

Any local behavior k induces a *global behavior* $k^\dagger : A^G \rightarrow A^G$, a map between configurations, via $k^\dagger(c)(x) =_{\text{df}} k(c \triangleright x)$ where $\triangleright : A^G \times G \rightarrow A^G$ (the *translation* of configurations) is defined by $(c \triangleright x)(y) =_{\text{df}} c(x \cdot y)$. The translation is a uniformly continuous function. It follows that the global behavior k^\dagger is also uniformly continuous.

Local behaviors on a fixed universe G and fixed alphabet A form a monoid with unit jd given by $\text{jd}(c) =_{\text{df}} c(1_G)$ and multiplication \bullet given by $\ell \bullet k =_{\text{df}} \ell \circ k^\dagger$. Indeed, it is easy to see that jd is uniformly continuous and \bullet preserves uniform continuity (because $(-)^{\dagger}$ does) and the monoid laws turn out to hold too.

We now make two small generalizations and make a richer category out of local behaviors: after all, a monoid is a category with one object. First, we do not insist that the alphabet be a discrete uniform space, it may be any uniform space. And second, we give up the idea of a fixed alphabet: we let the local behavior change the alphabet.

For a fixed monoid G (the universe), we redefine a *local behavior* between two general uniform spaces (the source and target alphabets) A and B to be a uniformly continuous function $k : A^G \rightarrow B$ where A^G is given the product uniformity.

Local behaviors now make a category that has as objects alphabets and as maps local behaviors between them. The identity on A is $\text{jd}_A : A^G \rightarrow A$ given by $\text{jd}_A(c) =_{\text{df}} c(1_G)$ and the composition $\ell \bullet k : A^G \rightarrow C$ of two maps $k : A^G \rightarrow B$

¹Instead of the monoid, one usually takes a group in the cellular automata literature. But we do not need inverses in this paper.

²The alphabet is often required to be finite. We make this assumption only where we need it.

³(Cf. [5, Th. 1.9.1]) Every entourage of the prodiscrete uniformity contains an entourage of the form $V_N = \{(c, e) \mid c|_N = e|_N\}$ with $N \subseteq G$ finite. If $k : A^G \rightarrow A$ is u.c. with A uniformly discrete, then $(k \times k)(V_N) \subseteq \Delta$ for some finite $N \subseteq G$: thus, $k(c)$ only depends on $c|_N$.

and $\ell : B^G \rightarrow C$ given by $\ell \bullet k =_{\text{df}} \ell \circ k^\dagger$ where $k^\dagger : A^G \rightarrow B^G$ is defined by $k^\dagger(c)(x) =_{\text{df}} k(c \triangleright_A x)$ from $\triangleright_A : A^G \times G \rightarrow A^G$ defined by $(c \triangleright_A x)(y) =_{\text{df}} c(x \cdot y)$. Notice that these definitions coincide exactly with those we made for the monoid of local behaviors above, except that local behaviors can now mediate between different alphabets that need not be uniformly discrete. The function jd_A is still uniformly continuous for any A and the operation \bullet preserves uniform continuity.

While the generalized definition of local behaviors is more liberal than the classical one, it is conservative over it in the following sense: The local behaviors from any uniformly discrete space A back to itself are exactly the classical local behaviors on A seen as a set.

We will now recover our category of local behaviors from a categorical generality, by showing that it is a straightforward instance of the coKleisli construction for a comonad.

Any fixed monoid $(G, 1_G, \cdot)$ determines a comonad (D, ε, δ) on **Unif** (in fact, on any category where the carrier G is exponentiable, so also, e.g., on **Set** and **Top**) (the *cellular automata* or *exponent comonad*) as follows. The object mapping part of D is defined by $DA =_{\text{df}} A^G$, where A^G is the G -exponential of A , i.e., the space of uniformly continuous functions from G to A equipped with the prodiscrete uniformity. The morphism mapping part is defined by $Df =_{\text{df}} f^G$, i.e., $Df(c) =_{\text{df}} f \circ c$. The components of the counit $\varepsilon_A : A^G \rightarrow A$ and comultiplication $\delta_A : A^G \rightarrow (A^G)^G$ are defined by $\varepsilon_A(c) =_{\text{df}} c(1_G)$ and $\delta_A(c)(x) =_{\text{df}} c \triangleright_A x$ (so that $\delta_A(c)(x)(y) = c(x \cdot y)$); these functions are uniformly continuous. The general definition of the coKleisli extension $(-)^{\dagger}$ via the morphism mapping part of D and comultiplication δ tells us that $k^\dagger(c)(x) = Dk(\delta_A(c))(x) = k(\delta_A(c)(x)) = k(c \triangleright_A x)$.

The laws of a comonad are proved from the monoid laws for G by the following calculations (we omit the proofs of the naturality conditions of ε and δ).

$$\begin{aligned} \varepsilon_{DA}(\delta_A(c))(x) &= \delta_A(c)(1_G)(x) = c(1_G \cdot x) = c(x) \\ c(x) &= c(x \cdot 1_G) = \delta_A(c)(x)(1_G) = \varepsilon_A(\delta_A(c)(x)) = D\varepsilon_A(\delta_A(c))(x) \\ \delta_{DA}(\delta_A(c))(x)(y)(z) &= \delta_A(c)(x \cdot y)(z) = c((x \cdot y) \cdot z) \\ &= c(x \cdot (y \cdot z)) = \delta_A(c)(x)(y \cdot z) = \delta_A(\delta_A(c)(x))(y)(z) = D\delta_A(\delta_A(c))(x)(y)(z) \end{aligned}$$

As we have seen, a comonad on a category always defines two canonical splittings of its underlying functor into two adjoint functors. The coKleisli splitting of our comonad D on **Unif** goes via the coKleisli category which has as objects those of **Unif** and as maps from A to B those from DA to B in **Unif**. The identity on A is $\text{jd}_A =_{\text{df}} \varepsilon_A$ and the composition of k and ℓ is $\ell \bullet k =_{\text{df}} \ell \circ k^\dagger$. Note that these are exactly the data of the category of local behaviors that we introduced above. But this time we do not have to prove that the unital and associativity laws of the category hold. Our proof obligations went into establishing that the comonad data are well defined and the comonad laws hold.

5. Retrieving the Curtis-Hedlund theorem

Let (D, ε, δ) be the G -exponential comonad on **Unif** for a given monoid $(G, 1_G, \cdot)$, with G endowed with the discrete uniformity, as introduced in the previous section.

As we know from Section 2, $\text{coKl}(D)$ is equivalent to the category of cofree D -coalgebras under a comparison functor E that sends a coKleisli map (local behavior) $k : DA \rightarrow B$ to the cofree coalgebra map $k^\dagger : (DA, \delta_A) \rightarrow (DB, \delta_B)$, which, as a

map of **Unif**, we know to be the corresponding global behavior. Hence, a map $f : DA \rightarrow DB$ would be a global behavior if and only if f is a cofree coalgebra map.

Now, given an arbitrary comonad, it is usually of interest to study its general coalgebras and not only the cofree ones. We too follow this thumb rule.

By definition, objects in $\text{coEM}(D)$ are pairs of objects A and maps $u : A \rightarrow DA$ in **Unif** satisfying

$$\begin{array}{ccc} A & \xrightarrow{u} & A^G \\ & \searrow & \downarrow \varepsilon_A \\ & & A \end{array} \quad \begin{array}{ccc} A & \xrightarrow{u} & A^G \\ u \downarrow & & \downarrow \delta_A \\ A^G & \xrightarrow{u^G} & (A^G)^G \end{array}$$

In our case, the first equation simply means $u(a)(1_G) = a$ while the second one simplifies to $u(a)(x \cdot y) = u(u(a)(x))(y)$. This writing, however, is cumbersome and unexplicative.

To see more, we uncurry $u : A \rightarrow A^G$ to $\otimes : A \times G \rightarrow A$, so that $a \otimes x = u(a)(x)$. Then the two equations become $a \otimes 1_G = a$, and $a \otimes (x \cdot y) = (a \otimes x) \otimes y$. Diagrammatically, this is to require commutation of

$$\begin{array}{ccc} A & \xrightarrow{\rho_A} & A \times 1 \xrightarrow{A \times 1_G} & A \times G \\ & \searrow & & \downarrow \otimes \\ & & & A \end{array} \quad \begin{array}{ccc} (A \times G) \times G & \xrightarrow{\alpha_{A,G,G}} & A \times (G \times G) \xrightarrow{A \times (\cdot)} & A \times G \\ \otimes \times G \downarrow & & & \downarrow \otimes \\ A \times G & \xrightarrow{\otimes} & & A \end{array}$$

where ρ and α are the right unital and associative laws of the product monoidal structure. But these are precisely the laws of a (*right*) action of G on A (where A is a uniform space, so we expect an action to also be uniformly continuous).

Let us now consider coalgebra maps. A map $f : (A, u) \rightarrow (B, v)$ in $\text{coEM}(D)$ is a map in **Unif** that commutes with u and v or (which is equivalent) with their uncurried forms \otimes and \circlearrowright as shown on the left and right diagrams below, respectively:

$$\begin{array}{ccc} A & \xrightarrow{u} & A^G \\ f \downarrow & & \downarrow f^G \\ B & \xrightarrow{v} & B^G \end{array} \quad \begin{array}{ccc} A \times G & \xrightarrow{\otimes} & A \\ f \times G \downarrow & & \downarrow f \\ B \times G & \xrightarrow{\circlearrowright} & B \end{array}$$

Clearly, coalgebra maps are just action maps.

We are now ready to consider maps of cofree D -coalgebras. The uncurried form of $\delta_A : A^G \rightarrow (A^G)^G$ is $\triangleright_A : A^G \times G \rightarrow A^G$. By what we have shown, a map $f : A^G \rightarrow B^G$ in **Unif** is a map between (A^G, δ_A) and (B^G, δ_B) if and only if $f(c \triangleright_A x) = (f(c) \triangleright_B x)$. We see that maps between the cofree coalgebras (A^G, δ_A) and (B^G, δ_B) are precisely those maps $f : A^G \rightarrow B^G$ that commute with the translation!

With our reasoning, we have reproved the version of Curtis-Hedlund theorem given by Ceccherini-Silberstein and Coornaert ([4, Th. 1.1], [5, Th. 1.9.1]): global behaviors between uniform spaces A and B are those uniformly continuous functions between the product uniformities on A^G and B^G that commute with the translation. Keep in mind that we are allowing arbitrary uniform spaces as alphabets, and speaking of global behaviors in the sense of Section 4. It is, however, immediate to specialize to the statement of C.-S. and C. by requiring A and B to be uniformly discrete. Then local behaviors are precisely the finitary functions from A^G to B .

The original Curtis-Hedlund theorem ([10], [5, Th. 1.8.1]) corresponds to the special case where A is finite and discrete. In this case, A^G is compact and any continuous function between A^G and B^G is then uniformly continuous. We conclude,

in this case, that global behaviors between A to B are those continuous functions between the product topologies on A^G and B^G that commute with the translation.

We have seen that the cofree coalgebra $(B^G, \delta_B : B^G \rightarrow (B^G)^G)$ on B is the currying of the translation action $(B^G, \triangleright_B : B^G \times G \rightarrow B^G)$. Because of the cofreeness on B of this coalgebra, the action must also be the cofree.

In basic terms, this means that for any action $(A, \otimes : A \times G \rightarrow A)$ and map $p : A \rightarrow B$, there is a unique action morphism f to the translation action such that $\varepsilon_B \circ f = p$, i.e., a map $f : A \rightarrow B^G$ such that the following diagrams commute:

$$\begin{array}{ccc} A \times G & \xrightarrow{\otimes} & A \\ f \times G \downarrow & & \downarrow f \\ B^G \times G & \xrightarrow{\triangleright_B} & B^G \xrightarrow{\varepsilon_B} B \end{array}$$

This fact can of course be proved from first principles, but we learned it for free!

6. Retrieving the reversibility principle

We will now recover the reversibility principle. Let again $(G, 1_G, \cdot)$ be a monoid and (D, ε, δ) be the G -exponential comonad on **Unif** as introduced in Section 4.

Suppose we have a cofree coalgebra map $f : (DA, \delta_A) \rightarrow (DB, \delta_B)$ so that f has an inverse f^{-1} as a map of **Unif**.

A very simple diagram chase (not specific to our particular comonad; we only use that D is a functor!) shows that f^{-1} is also a cofree coalgebra map, i.e., an inverse of f in $\text{coEM}(D)$:

$$\begin{array}{ccccc} & & DA & \xrightarrow{\delta_A} & D(DA) \equiv D(DA) \\ & \nearrow f^{-1} & \downarrow f & & \downarrow Df \\ DB & \equiv & DB & \xrightarrow{\delta_B} & D(DB) \xrightarrow{Df^{-1}} D(DA) \end{array}$$

We have thus reproved the reversibility principle [5, Th. 1.10.1]: A global behavior $f : A^G \rightarrow B^G$ (a uniformly continuous function commuting with the translation) between uniform spaces A, B is reversible if f has a uniformly continuous inverse.

If both A and B are finite and discrete, an inverse f^{-1} of f is necessarily uniformly continuous, because, in this case, A^G and B^G are compact Hausdorff and f^{-1} is continuous. So we obtain a special case [5, Th. 1.10.2]: A global behavior $f : A^G \rightarrow B^G$ (a continuous function commuting with the translation) between uniform spaces A, B is reversible if f has an inverse.

In general, an inverse of a uniformly continuous function is not necessarily uniformly continuous. For a counterexample, see [5, Example 1.10.3].

For reversibility, it is useful, if the monoid G is actually a group. In particular, for reversibility of $\delta_A : DA \rightarrow D(DA)$, G must be a group.

7. Distributive laws and 2-dimensional cellular automata

We will now proceed to two variations on the theme of cellular automata as co-Kleisli maps—2-dimensional (classical) cellular automata and point-dependent cellular automata. In both cases we first introduce some further comonad theory relevant for our cause.

Sometimes, but not always, the composition D^1D^0 of two comonads D^0 and D^1 on the same category \mathcal{C} is a comonad. It is the case, if there is a distributive law of D^1 over D^0 . A *distributive law* of a comonad $(D^1, \varepsilon^1, \delta^1)$ over a comonad $(D^0, \varepsilon^0, \delta^0)$ is a natural transformation $\kappa : D^1D^0 \rightarrow D^0D^1$ making the diagrams

$$\begin{array}{ccc}
D^1D^0 & \xrightarrow{\kappa} & D^0D^1 \\
D^1\varepsilon^0 \searrow & & \swarrow \varepsilon^0D^1 \\
& D^1 & \\
\end{array}
\qquad
\begin{array}{ccc}
D^1D^0 & \xrightarrow{\kappa} & D^0D^1 \\
D^1\delta^0 \downarrow & & \downarrow \delta^0D^1 \\
D^1D^0D^0 & \xrightarrow{\kappa D^0} & D^0D^1D^0 \xrightarrow{D^0\kappa} D^0D^0D^1
\end{array}$$

$$\begin{array}{ccc}
D^1D^0 & \xrightarrow{\kappa} & D^0D^1 \\
\varepsilon^1D^0 \searrow & & \swarrow D^0\varepsilon^1 \\
& D^0 & \\
\end{array}
\qquad
\begin{array}{ccc}
D^1D^0 & \xrightarrow{\kappa} & D^0D^1 \\
\delta^1D^0 \downarrow & & \downarrow D^0\delta^1 \\
D^1D^1D^0 & \xrightarrow{D^1\kappa} & D^1D^0D^1 \xrightarrow{\kappa D^1} D^0D^1D^1
\end{array}$$

commute. A distributive law induces a comonad (D, ε, δ) defined by $D =_{\text{df}} D^1D^0$, $\varepsilon =_{\text{df}} \varepsilon^1\varepsilon^0$, $\delta =_{\text{df}} D^1\kappa D^0 \circ \delta^1\delta^0$.

A distributive law also induces comonad liftings. For the lack of space, we concentrate on the coKleisli side of the picture. Given a distributive law κ of D^1 over D^0 , the comonad D^0 on \mathcal{C} lifts to $\text{coKl}(D^1)$, i.e., induces a comonad \bar{D}^0 on $\text{coKl}(D^1)$. This comonad is defined by: $\bar{D}^0A =_{\text{df}} D^0A$, $\bar{D}^0k =_{\text{df}} D^0k \circ \kappa_A : D^1D^0A \rightarrow D^0B$ (for $k : D^1A \rightarrow B$), $\varepsilon_A =_{\text{df}} \varepsilon_A^0 \circ \varepsilon_{D^0A}^1$, $\delta_A =_{\text{df}} \delta_A^0 \circ \varepsilon_{D^0A}^1$.

Via this lifting, the coKleisli category of the composite comonad D and the other ingredients of the coKleisli splitting of D can be obtained by a double coKleisli construction: We have $\text{coKl}(D) = \text{coKl}(\bar{D}^0)$ (equal strictly, not just isomorphic).

A simple example of a distributive law is obtained by taking D^1 to be any comonad and D^0 the product comonad defined by $D^0A =_{\text{df}} A \times C$ for a fixed object $C \in |\mathcal{C}|$. The distributive law κ of D^1 over D^0 is given by $\kappa_A =_{\text{df}} \langle D^1\pi_0, \varepsilon^1 \circ D^1\pi_1 \rangle : D^1(A \times C) \rightarrow D^1A \times C$. It follows that the functor D defined by $DA =_{\text{df}} D^1(A \times C)$ is a comonad.

Given now two monoids G_0, G_1 , we can think of a map $k : (A^{G_0})^{G_1} \rightarrow B$ in **Unif** as a “2-dimensional” (2D) cellular automaton on the universes G_0, G_1 between alphabets A and B (relying on the isomorphism $A^{G_0 \times G_1} \cong (A^{G_0})^{G_1}$).

Such a cellular automaton is by definition the same thing as a “1-dimensional” (1D) cellular automaton on the universe G_1 between alphabets A^{G_0} and B . Note that we can only see 2D cellular automata as 1D in this way, if we allow source and target alphabets of a cellular automaton to differ and if we do not require them to be uniformly discrete (notice that A^{G_0} carries the prodiscrete uniformity). But this view of 2D cellular automata as 1D, although nice, suffers from a serious drawback. Since the 1D views do not have the same source alphabets as the 2D originals, they do not compose the same way.

Distributive laws come to help. Let $D^0A =_{\text{df}} A^{G_0}$ and $D^1A =_{\text{df}} A^{G_1}$. There is a distributive law $\kappa : D^1D^0 \rightarrow D^0D^1$ defined by $\kappa_A(c)(x_1)(x_0) = c(x_0)(x_1) : (A^{G_0})^{G_1} \rightarrow (A^{G_1})^{G_0}$. Hence, the functor $DA =_{\text{df}} (A^{G_0})^{G_1}$ is a comonad, which is hardly a surprise. But there is more: We know that $\text{coKl}(D) = \text{coKl}(\bar{D}^0)$. Hence, a good view of $k : (A^{G_0})^{G_1} \rightarrow B$ as a 1D cellular automaton is not as a **Unif**-cellular automaton on the universe G_1 between the alphabets A^{G_0} and B , but on as a $\text{coKl}(D^1)$ -cellular automaton on the universe G_0 between the alphabets A and B . Then 1D views compose exactly as their 2D originals. We see that it makes sense to consider maps of categories other than **Unif**!

1D views of 2D cellular automata were of interest to Dennuzio et al. [6]

8. Comonad maps and point-dependent cellular automata

To make a category out of comonads over a fixed category \mathcal{C} one needs a suitable notion of comonad maps. A *comonad map* between two comonads (D, ε, δ) and $(D', \varepsilon', \delta')$ on \mathcal{C} is a natural transformation $\tau : D \rightarrow D'$ making the diagrams

$$\begin{array}{ccc} D & \xrightarrow{\varepsilon} & \text{Id}_{\mathcal{C}} \\ \tau \downarrow & & \nearrow \\ D' & \xrightarrow{\varepsilon'} & \end{array} \quad \begin{array}{ccc} D & \xrightarrow{\delta} & DD \\ \tau \downarrow & & \downarrow \tau\tau \\ D' & \xrightarrow{\delta'} & D'D' \end{array}$$

commute. Comonads and comonad maps on \mathcal{C} form a category. The identity and composition of comonad maps is inherited from the category of natural transformations between endofunctors on \mathcal{C} .

A comonad map τ between (D, ε, δ) and $(D', \varepsilon', \delta')$ relates the coKleisli and coEilenberg-Moore categories between the two comonads. It defines a functor from $\text{coKl}(D')$ to $\text{coKl}(D)$ and a functor from $\text{coEM}(D)$ to $\text{coEM}(D')$.

We now introduce point-dependent cellular automata (studied under the name of non-uniform cellular automata by Cattaneo et al. [3]). For a set G , the local behavior of a point-dependent cellular automaton between uniform spaces A, B is a uniformly continuous function $k : A^G \times G \rightarrow B$. Note the added second argument compared to the definition of a classical local behavior.

It turns out that local and global behaviors of point-dependent cellular automata can be analyzed in the same way as those of classical cellular automata. In particular, their local behaviors are the same thing as coKleisli maps of a suitable comonad (D, ε, δ) on \mathbf{Unif} and global behaviors are the corresponding cofree coalgebra maps!

Let us review the data of the comonad. The object mapping of D is defined by $DA =_{\text{df}} A^G \times G$ and the morphism mapping by $Df =_{\text{df}} f^G \times G$, i.e., $Df(c, x) =_{\text{df}} (\lambda y. f(c(y)), x)$. The components of the counit and comultiplication $\varepsilon_A : A^G \times G \rightarrow A$ and $\delta_A : A^G \times G \rightarrow (A^G \times G)^G \times G$ are defined by $\varepsilon_A(c, x) =_{\text{df}} c(x)$, $\delta_A(c, x) = (\lambda y. (c, y), x)$. Accordingly, the coKleisli extension $k^\dagger : A^G \times G \rightarrow B^G \times G$ of a map $k : A^G \times G \rightarrow B$ is forced to satisfy $k^\dagger(c, x) = Dk(\delta(c, x)) = Dk(\lambda y. (c, y), x) = (\lambda y. k(c, y), x)$.

When is a map $f : A^G \times G \rightarrow B^G \times G$ a global behavior? It is a global behavior iff it is a cofree coalgebra map. Not surprisingly at all, the conditions for f being a cofree coalgebra map reduce to the condition that $f(c, x) = (g(c), x)$ for some $g : A^G \rightarrow B^G$.

Assume G is endowed with a monoid structure $(1_G, \cdot)$. Let $(D', \varepsilon', \delta')$ be the comonad of classical cellular automata. The translation \triangleright is a comonad map from D to D' . Accordingly, any classical local behavior is also a point-dependent local behavior that simply makes no use the point information that is available.

9. Conclusions

It was not the purpose of this paper to prove deep or difficult theorems. Rather, we set out to experiment with definitions. We deem that this experiment succeeded. We were pleased to learn that, from the category-theoretic point-of-view, cellular automata are a “natural” construction with “natural” properties. Crucially, classical cellular automata are coKleisli maps of the exponential comonad on \mathbf{Unif} , and it is harmless to accept alphabets with nondiscrete uniformities and variation

of alphabets, once it has been decided that local behaviors are uniformly continuous functions. But other base categories can be useful too, as the example of 2-dimensional cellular automata as 1-dimensional shows.

We hope to be able to extend this work to cover more results of cellular automata theory, in particular results toward the Garden of Eden theorem.

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